



PLAXIDITYX  
FORMERLY ARGUS

# Tales from a Penetration Testing Team

Insights from zero-day automotive vulnerabilities discovered in recent years



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# Introduction

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## In This Talk

- Found Pre-SOP & already fixed by the responsible parties
- Some of the details were **redacted** to ensure confidentiality of the involved parties



# Vulnerability #1:

Arbitrary Remote Code Execution over CAN

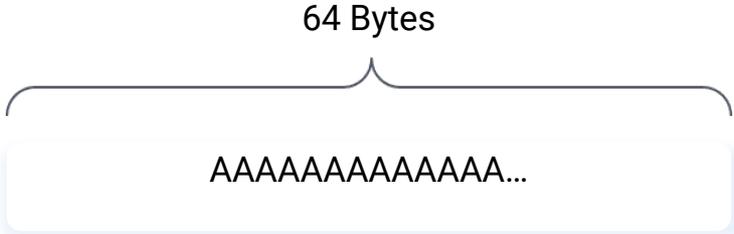
# Vulnerability #1

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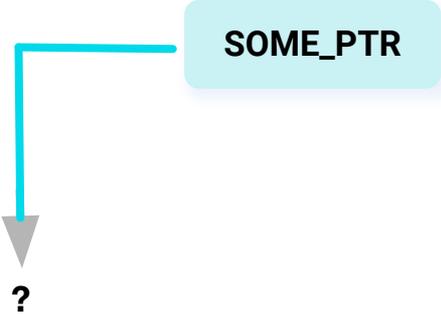
memcpy(**SOME\_PTR**, CANFD\_frame, 64)

# Vulnerability #1

```
memcpy(SOME_PTR, CANFD_frame, 64)
```

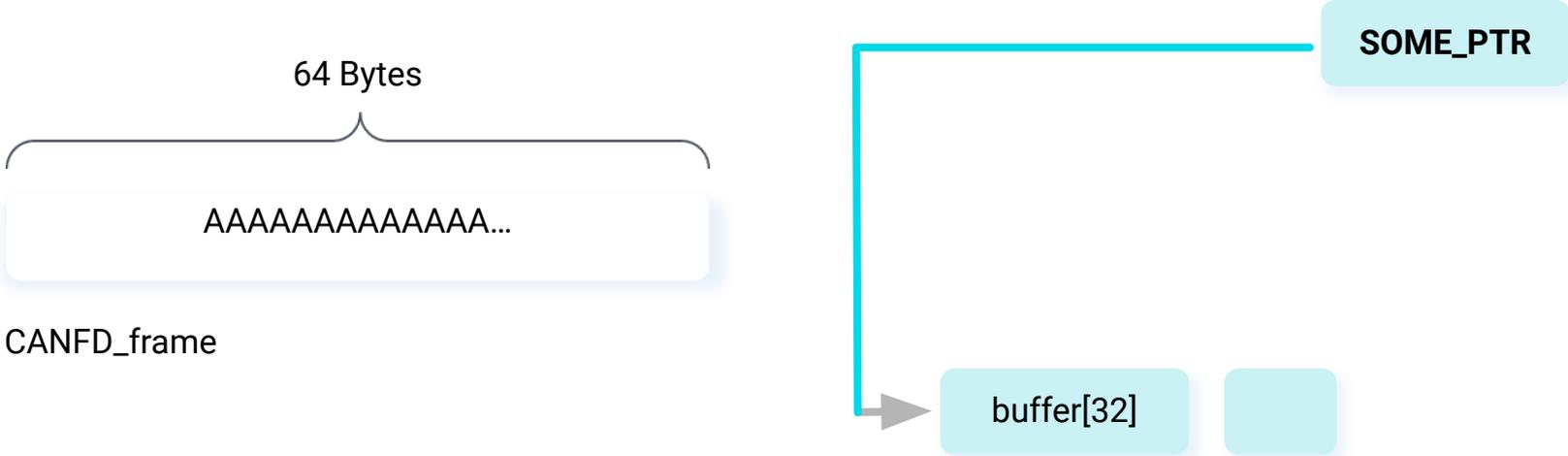


CANFD\_frame



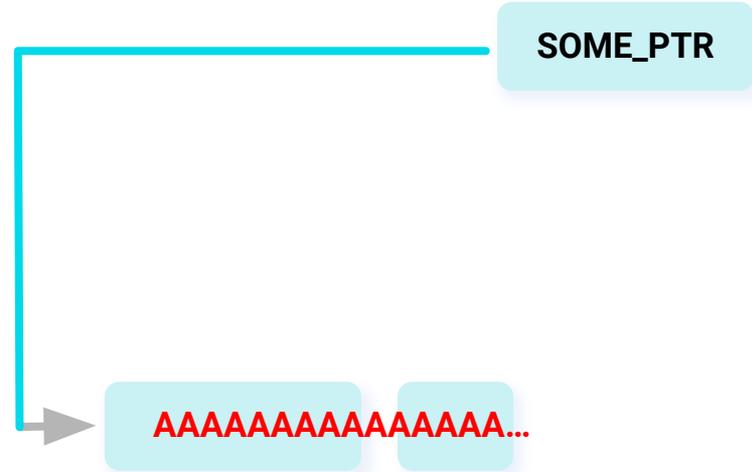
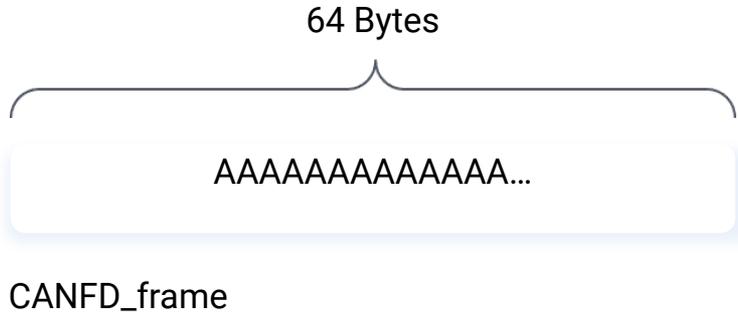
# Vulnerability #1

```
memcpy(SOME_PTR, CANFD_frame, 64)
```



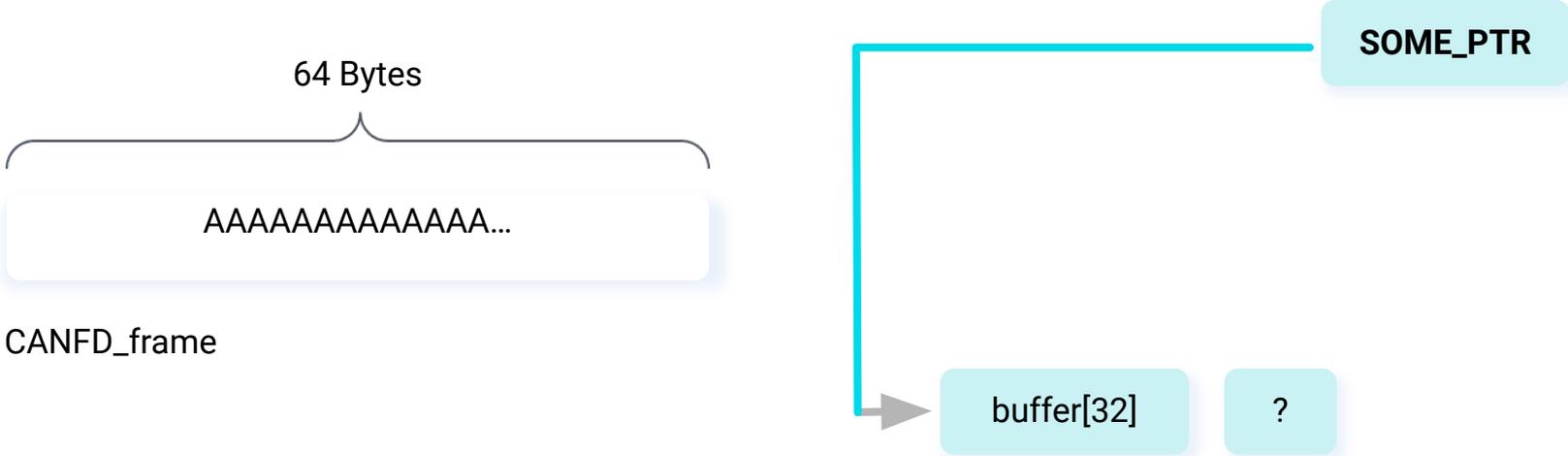
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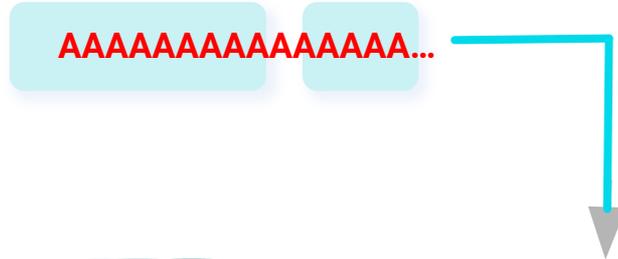
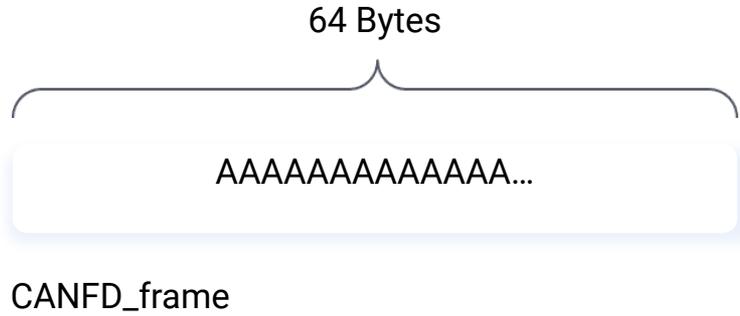
# Vulnerability #1

```
memcpy(SOME_PTR, CANFD_frame, 64)
```



# Vulnerability #1

```
memcpy(AAAAAAAA, CANFD_frame, 64)
```



# Vulnerability #1

```
ISR(){
```

```
...  
    memcpy(SOME_PTR, CANFD_frame, 64)  
...
```



# Vulnerability #1

---

```
ISR(){
```

```
...
```

```
    memcpy(SOME_PTR, CANFD_frame, 64)
```

```
...
```



Interrupt Service Routine (ISR) -> **Highest Privileges**

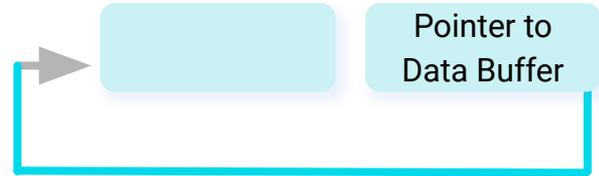
# Step 1

## Controlling the destination pointer

CAN-FD  
Frame #1

AAAAAAAAAAAAAAAACAFECAFE

Data  
Buffer



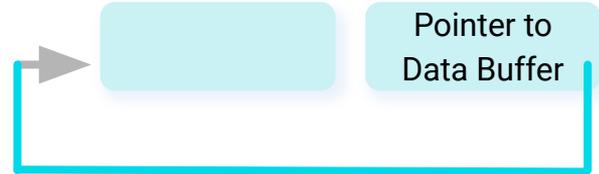
# Step 1

## Controlling the destination pointer

CAN-FD  
Frame #1

AAAAAAAAAAAAAAAACAFECAFE

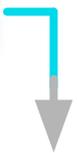
Data  
Buffer



Data  
Buffer

AAAAAAAAAAAAAAAA

CAFECAFE



# Step 2

## Writing

CAN-FD  
Frame #2

DEADBEEF

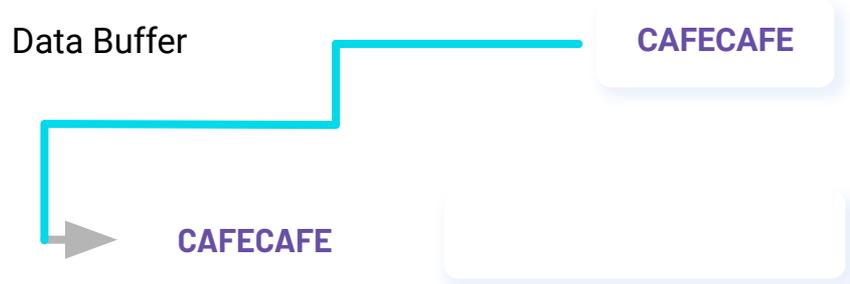
Data  
Buffer

CAFECAFE

CAFECAFE

# Step 2

## Writing



# Note

CAN-FD  
Frame #3

12 34 56 78

Data  
Buffer

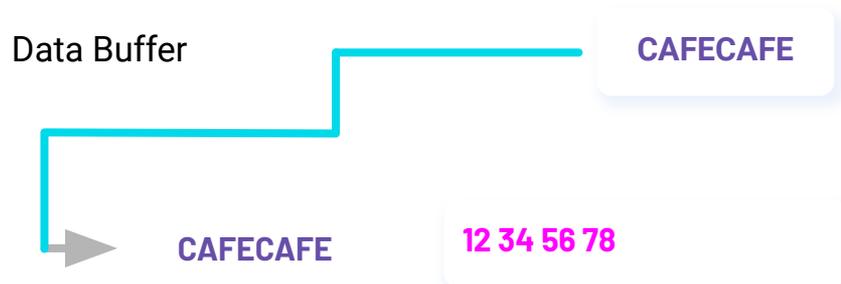
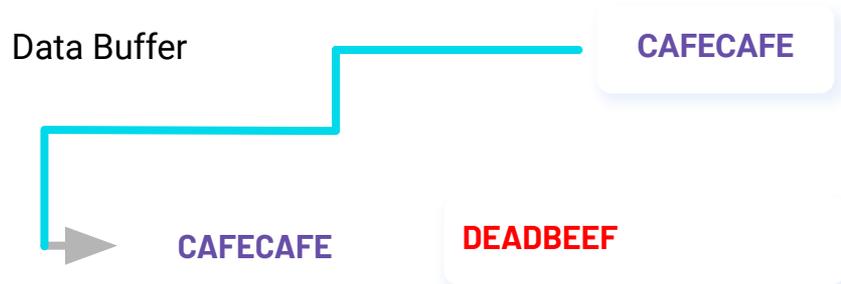
CAFECAFE

CAFECAFE

DEADBEEF

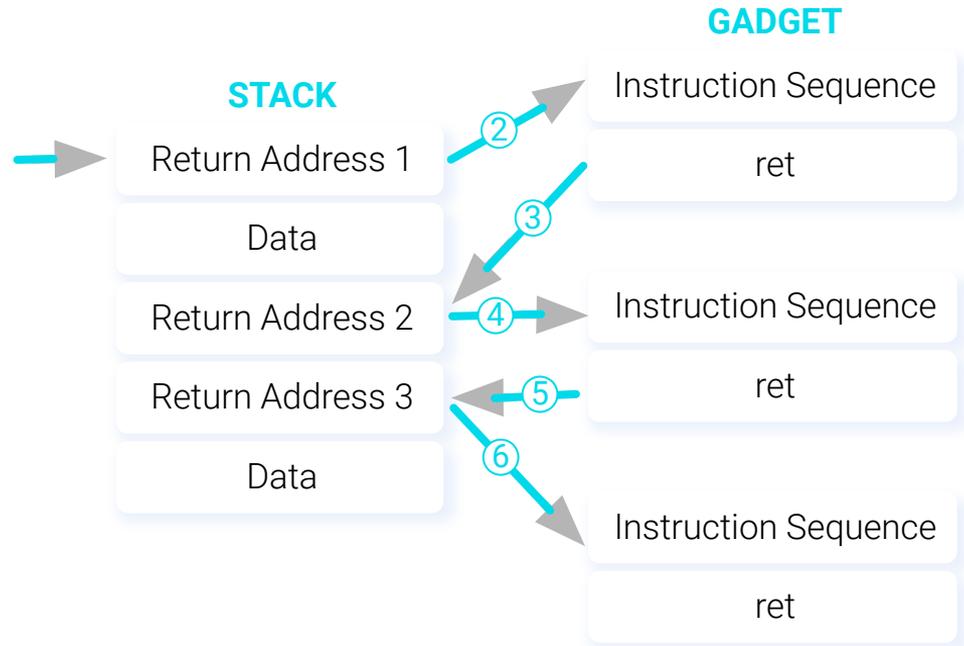
# Note

## Writing once



# Vulnerability #1 - Conclusion

- Arbitrary Write via CAN-FD messages
- Writing directly to the return address on the stack
- Bypassing stack canaries
- Use Return-Oriented-Programming to get the desired code flow



Springer: ROP-Hunt: Detecting Return-Oriented Programming Attacks in Applications

# Vulnerability #2:

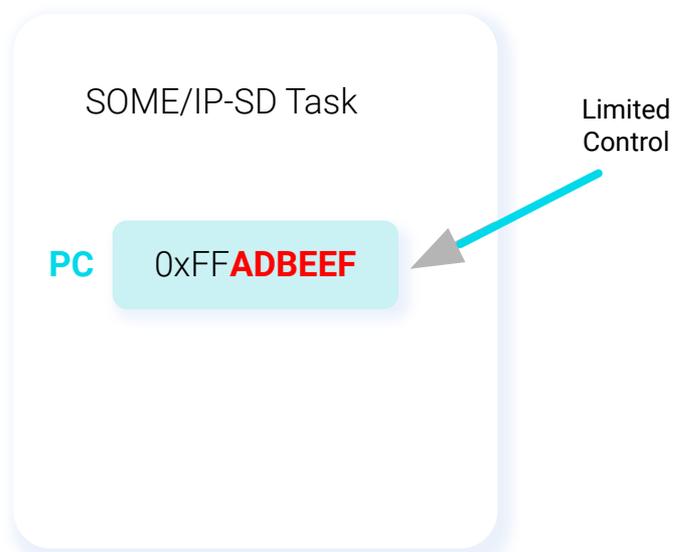
RCE over IPsec and SOME/IP-SD

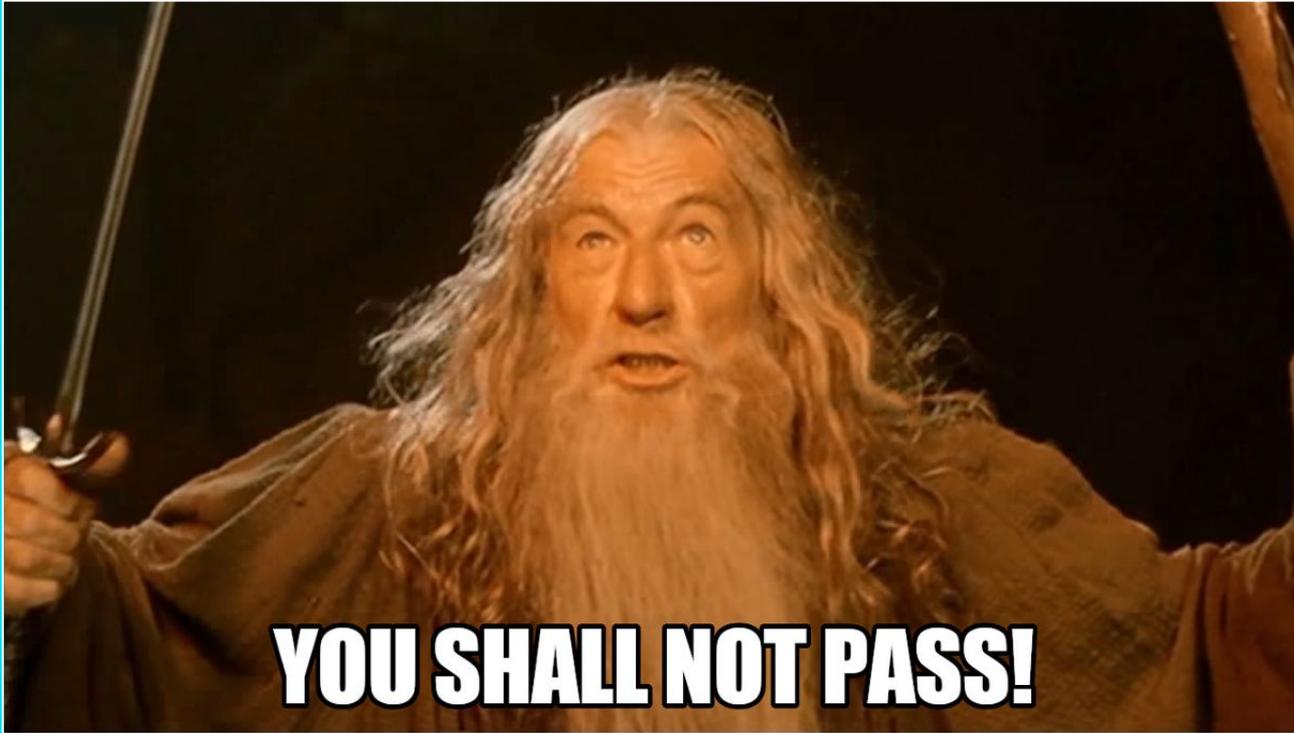
# Fuzzing SOME/IP-SD

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- We used our interface fuzzing tool to find issues in the SOME/IP-SD protocol handling
- The fuzzer successfully crashed the system
- A stack overflow caused the program counter to point to an invalid memory area
- No stack canaries were added as a countermeasure in the vulnerable function

# Vulnerability #2 - The SD task Problem



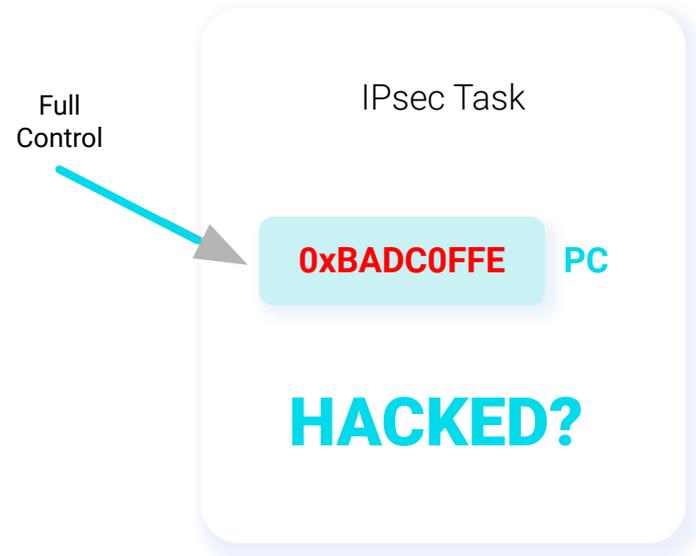


# Fuzzing IPsec

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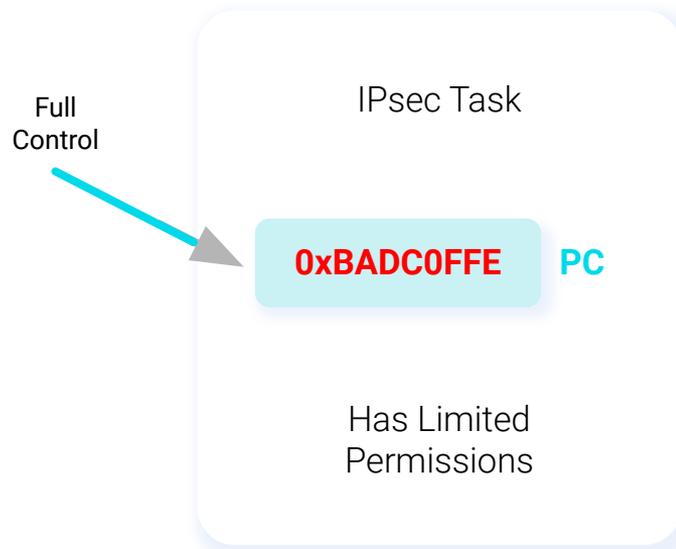
- IPsec was also part of the TCP/IP communication stack
- We used our interface fuzzing tool again
- Once more, we found a stack overflow that caused the program counter to point to an invalid memory area
- This happens as part of the authentication process (**No need to be authenticated!**)
- No stack canaries in the vulnerable function, again!

# Vulnerability #2 - Full Control

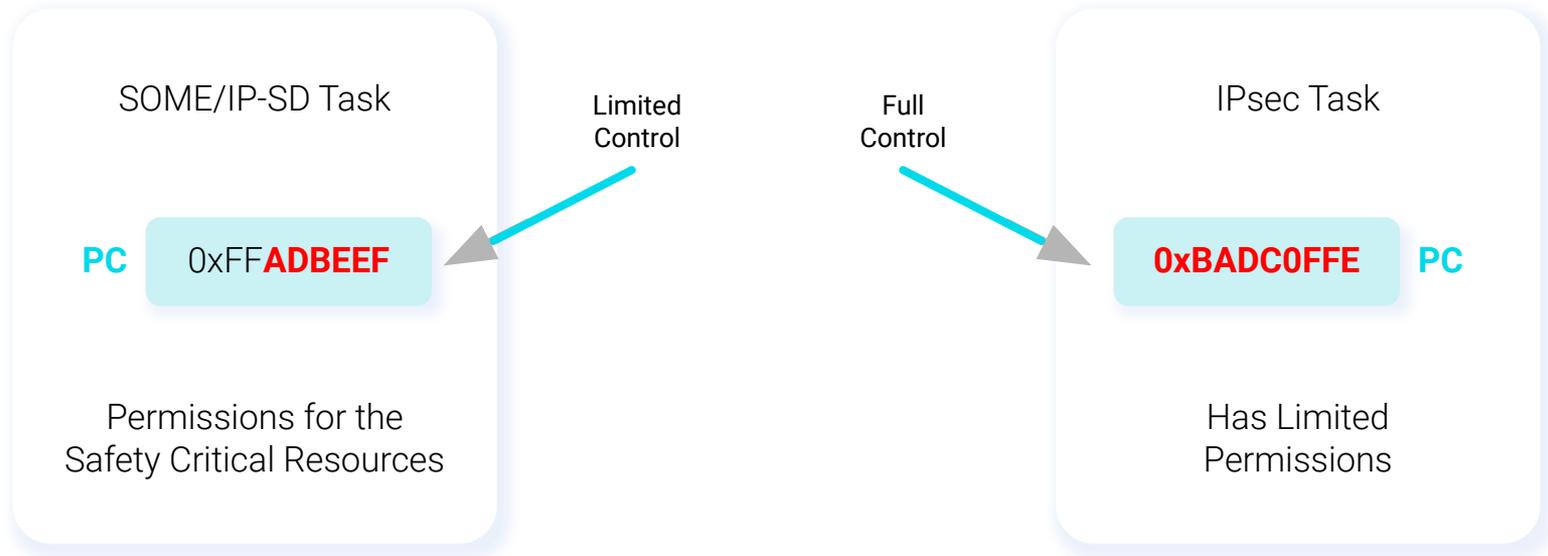




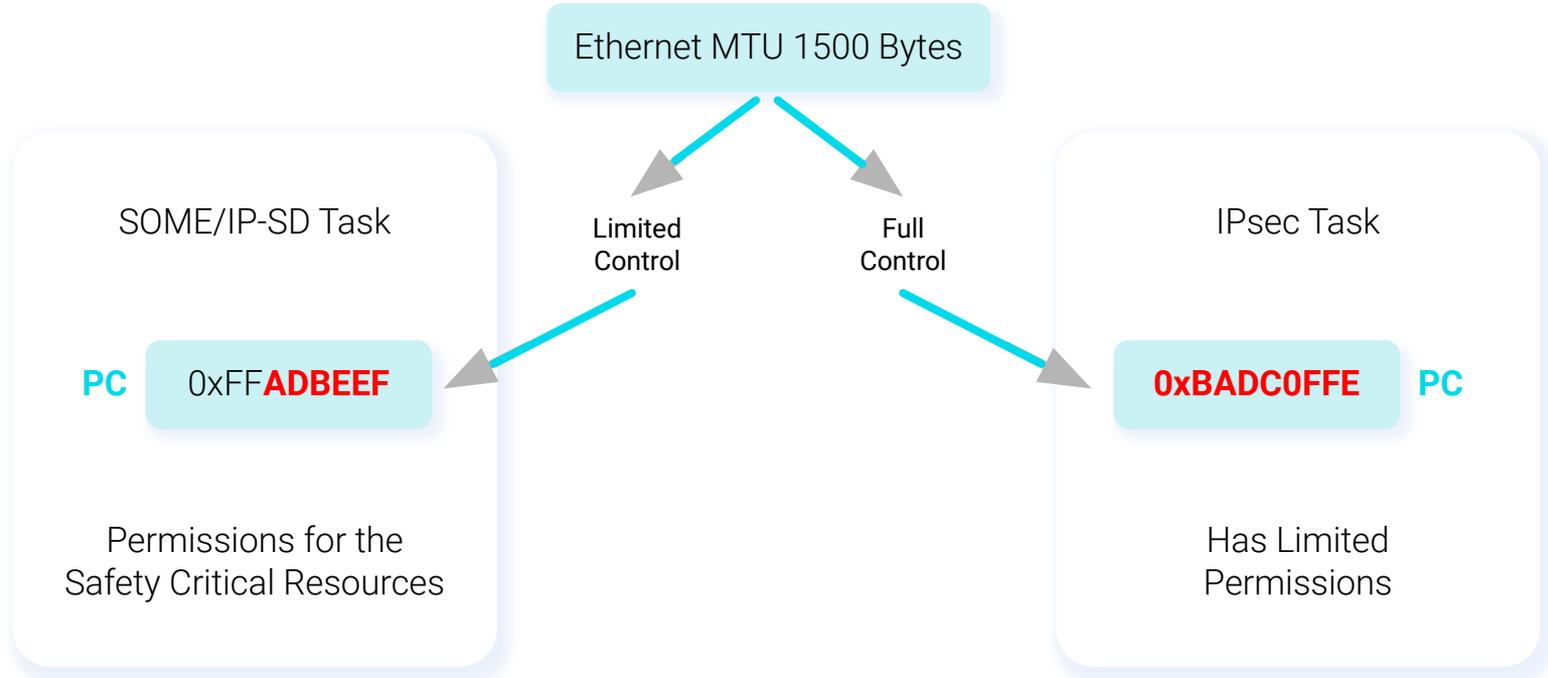
# Vulnerability #2 - The IPsec Task Problem



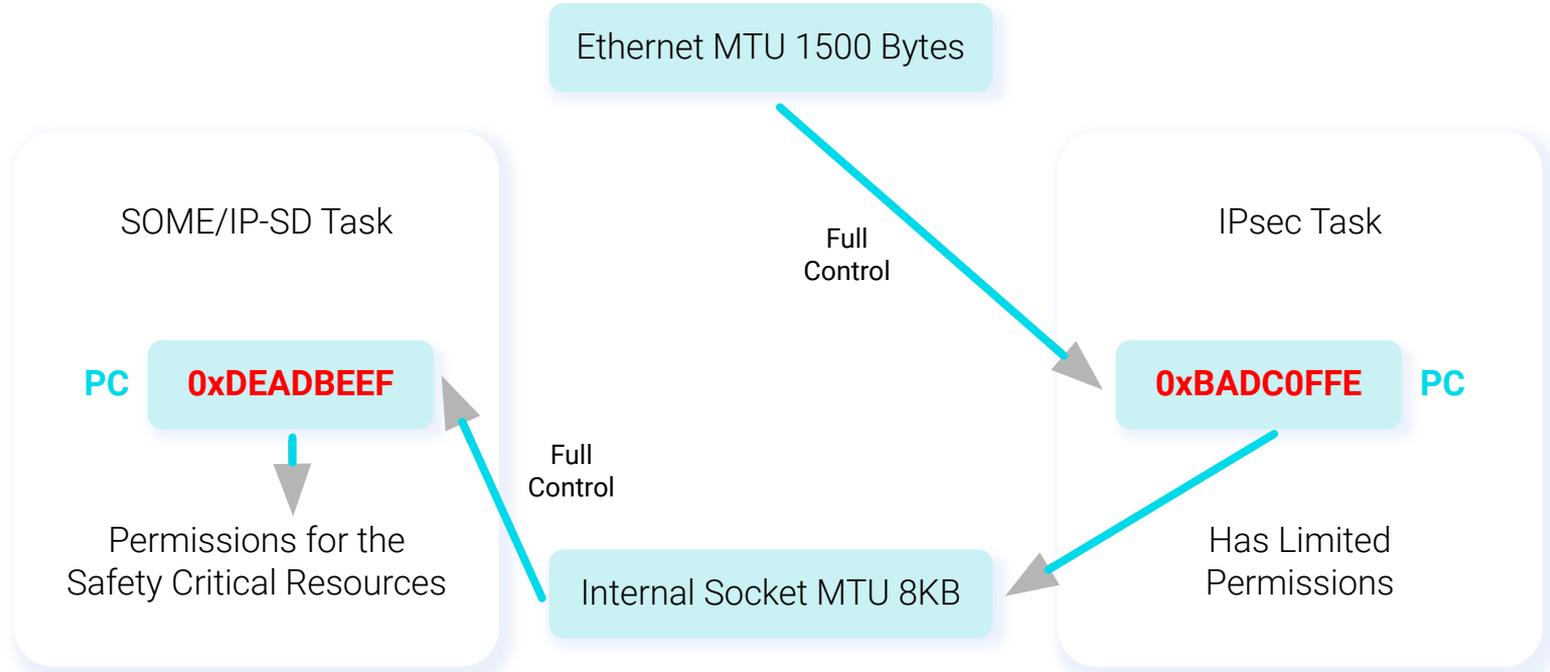
# Vulnerability #2 - The Problem



# Vulnerability #2 - The Problem



# Vulnerability #2 - Chaining Vulnerabilities





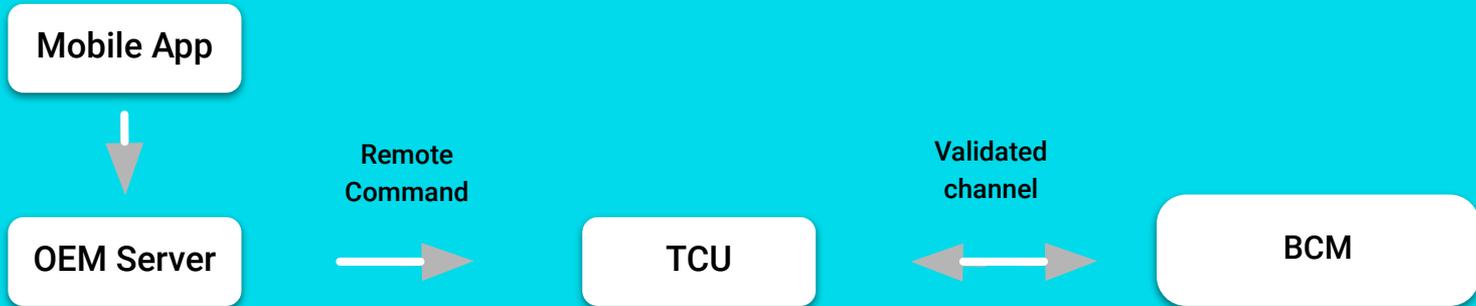
# AND THE UNION OF THE TWO TASKS

# Vulnerability #3:

## Shaky Cryptography

# Vulnerability #3 - Authentication of Remote Commands

- TCU can issue 'instructions' to the BCM based on received 'remote commands'
- BCM validates instructions through a unique pre-shared key and AES in order to prove they came from a valid TCU source.



# Typical Cryptography Related Gaps

## **Non Random -**

"Random" Values

## **Derivable Keys**

From Communication

## **Get / Set Key**

Functions

## **Hard Coded**

Master Key

## **Easily**

Disabled

## **Brute - Forceable**

Keys

# What Have We Found In This Setup

---

## Non Random -

"Random" Values

## Derivable Keys

From Communication



## Get / Set Key

Functions

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# What Have We Found In This Setup

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# What Have We Found In This Setup

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## Derivable Keys

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## Easily

Disabled

## Get / Set Key

Functions

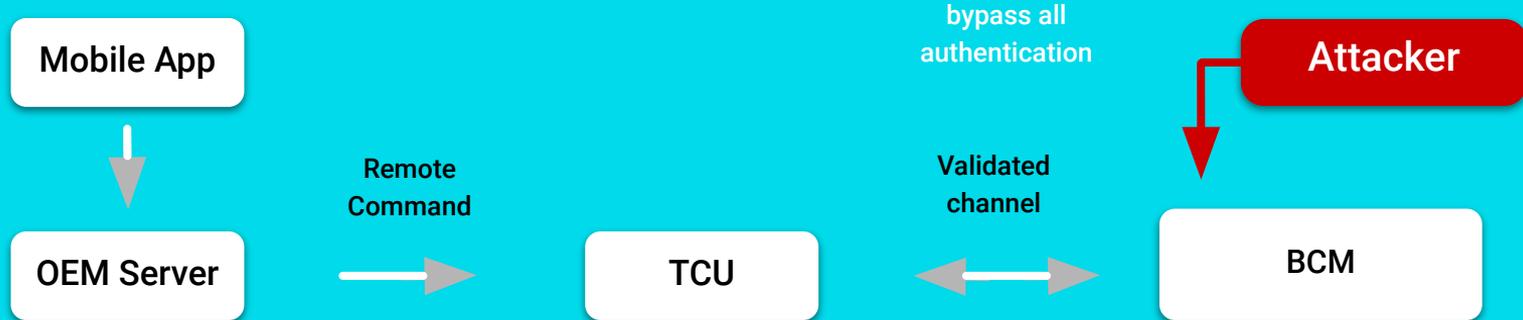


## Brute - Forceable

Keys

# Vulnerability #3 - Authentication of Remote Commands

- Using any of the above vulnerabilities we were able to bypass the validation and impersonate the TCU, without any prior knowledge of any key material

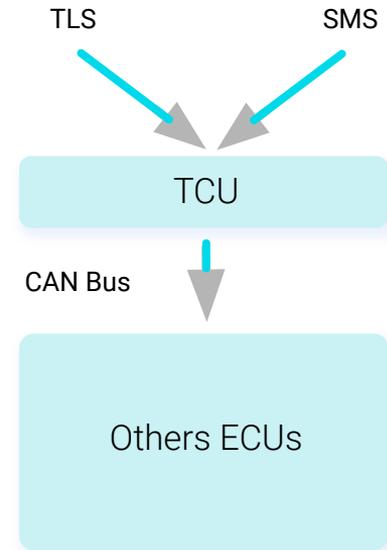


# Vulnerability #4:

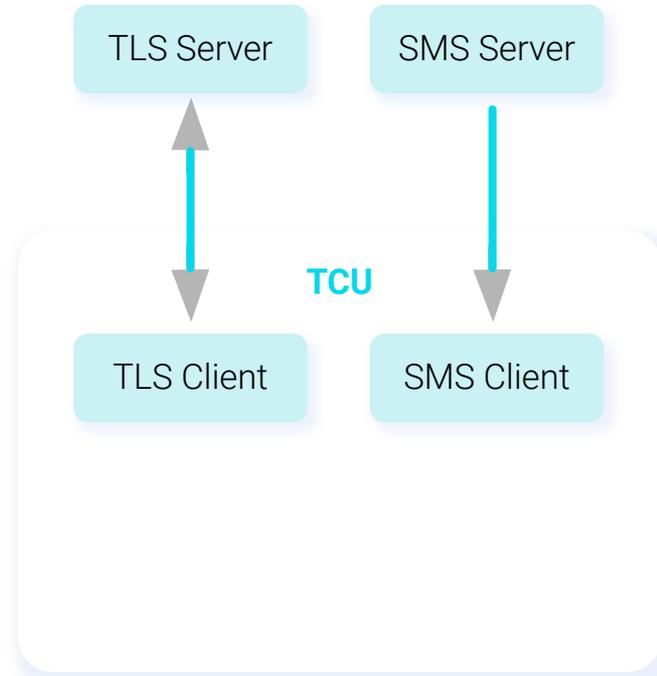
Remote and Persistent Vulnerability  
via Cellular Connection

# Vulnerability #4 - Setup

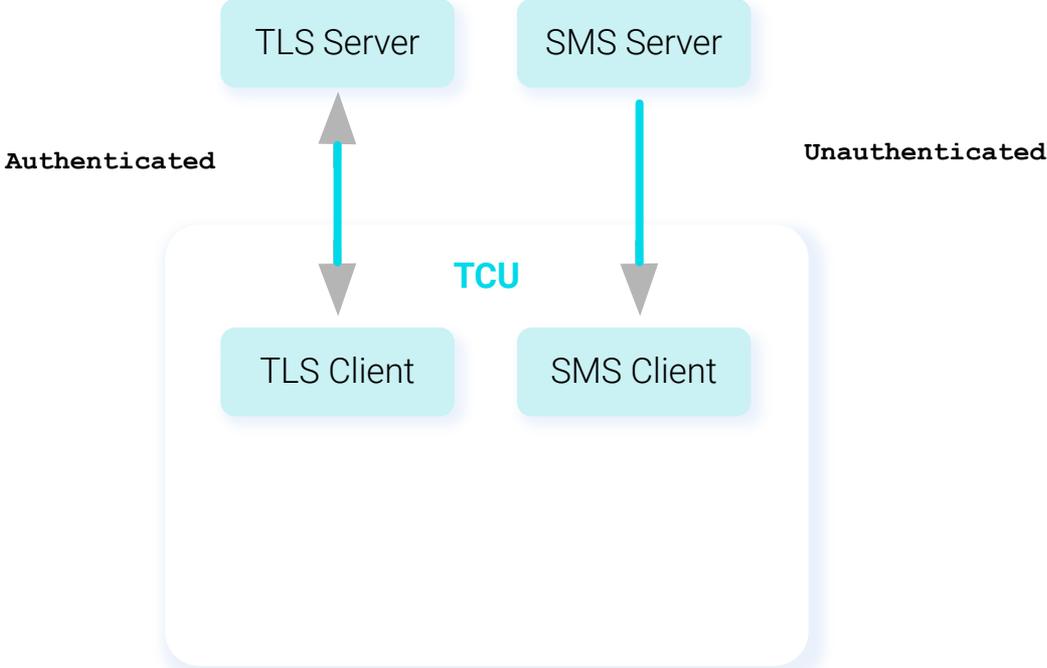
- TCU used in a heavy duty vehicle with a cellular modem
- Communication with the backend via secured TLS
- Connected directly to safety critical CAN bus
- Can get binary SMS to trigger communication with the backend



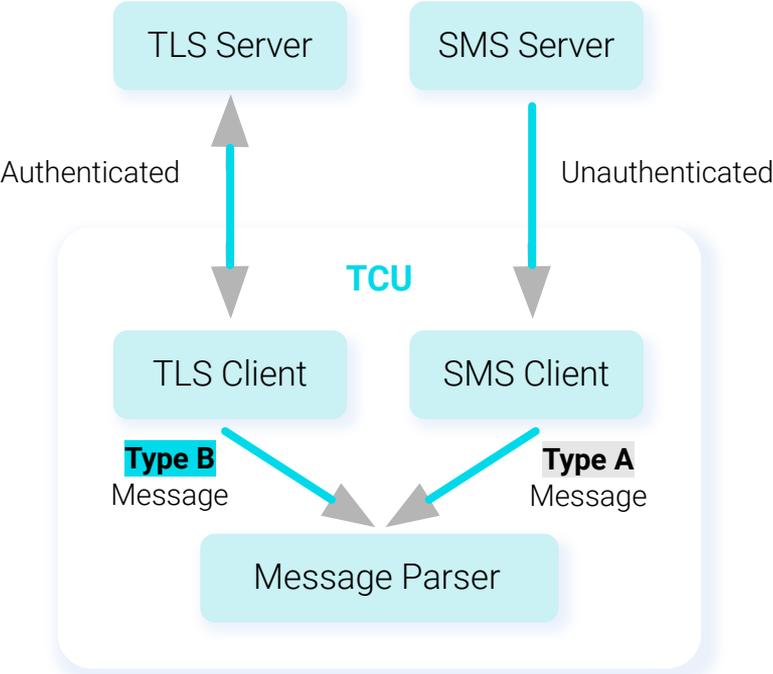
# Vulnerability #4 - Setup



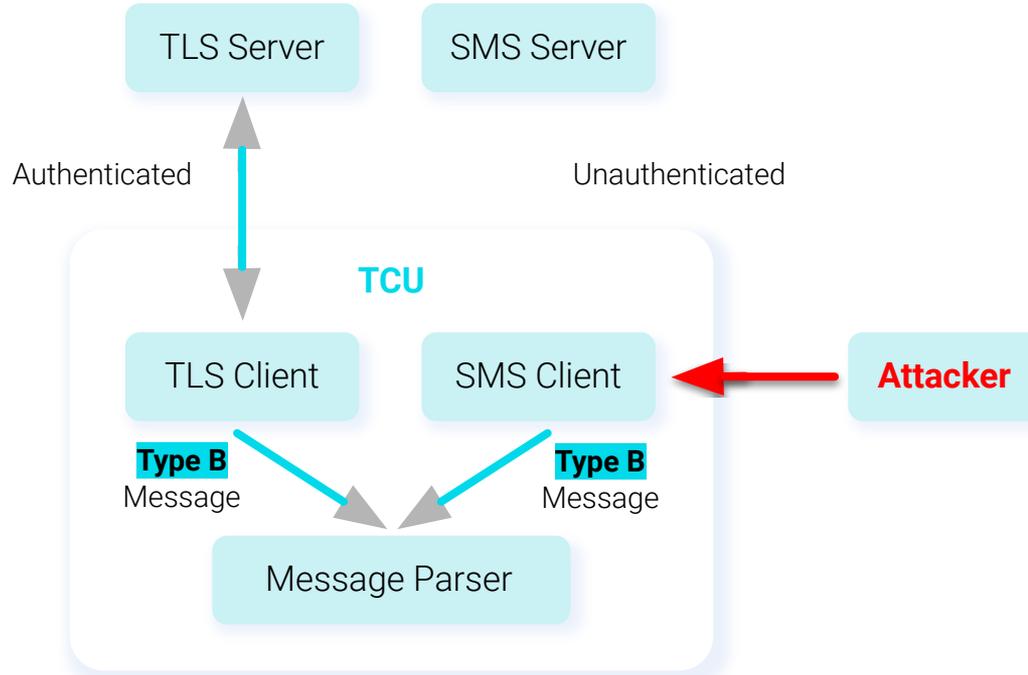
# Vulnerability #4 - Connections



# Vulnerability #4 - Messages Types



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# Vulnerability #4

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# Vulnerability #4

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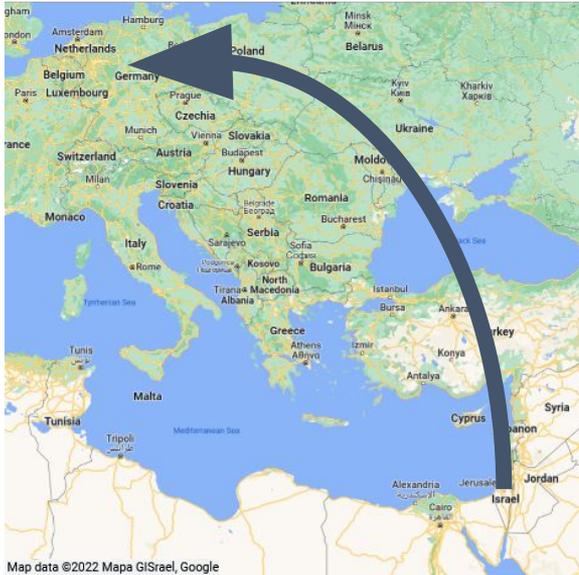
SMS

TLS

# Vulnerability #4 - The Next Step

- Attacker can use a **backend command to replace** a file within the system
- Some files on the system are **tagged with a digital-signature**
- But some of the partitions in the file system **did not have any validation**
- A subset of those were not flagged as **“noexec”**
- Therefore, an attacker **can inject an arbitrary executable** to any of these partitions

# Vulnerability #4 - End Game



- This vulnerability enabled us to run a full demo of injecting CAN messages to impact the vehicle via an unauthenticated binary SMS message. From **Tel Aviv, Israel** to **Europe**.

# Conclusions

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- High severity **Zero-Day vulnerabilities are still common in multiple types** of ECUs (some safety-critical)
- More **vulnerabilities are caught Pre-SOP** nowadays
- These vulnerabilities are the **result of faulty implementations of protocols** and interfaces
- Security of modern, **complex vehicle systems requires a holistic approach** with layered security controls and vast security auditing

# Conclusions

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**A**

Augmented security awareness during the development process

**A**

Fuzzing Tools

**A**

Don't underestimate a limited security issue

**A**

Hardening!

**A**

Penetration testing,  
by security researchers

# Thank you